

Objects in Scheme

We can use closures and mutation to do OOP in Scheme!

Recall the counter example from last class:

```
(define (make-counter)
 (let ((count 0))
   (lambda ()
      (set! count (+ 1 count))
      (display count)
      (newline))))
```

To add methods, the (first) argument to the returned lambda will be the *name* of the method we want to apply.

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More sophisticated counter:

```
(define (make-counter-obj)
 (let ((count 0))
  (lambda (command)
      (cond [(symbol=? command 'get) count]
        [(symbol=? command 'inc)
        (set! count (+ 1 count))]
        [(symbol=? command 'reset)
        (set! count 0)]))))
```

The object now has three methods: get, inc, and reset.

Built-in Data Structures Scheme has some useful built-in data structures: • Arrays (called "vectors"). (define A (make-vector 5)) (vector-set! A 3 'something) (vector-ref A 3) ; produces 'something (vector-ref A 5) ; error: out of bounds Hash tables (define H (make-hash)) (hash-set! H 2 'something) (hash-set! H (list 20 #f) 'crazy!) (hash-ref H '(20 #f)) ; produces 'crazy! (hash-ref H '(bad)) ; error: no key (bad) Fall 2011 4 / 11 Roche (USNA) SI413 - Class 6

Efficiency in Scheme: Fibonacci numbers Recall the problem of computing Fibonacci numbers from lab 1. (define (fib n) (if (<= n 1) 1 (+ (fib (- n 1)) (fib (- n 2))))) Why is this function so slow?

Memoization in Scheme Memoization is remembering the results of previous function calls. Why is functional programming *perfect* for memoization? In Scheme we can use vectors or hashes to provide this functionality. Some languages have built-in memoization.

Memoizing Fibonacci

Here's how we might memoize the Fibonacci function:

Stack space in recursive calls Recursive calls can use a lot of memory, even when the results are puny. ;; Sum of squares from 1 to n (define (ssq n) (if (= n 0) 0 (+ (sqr n) (ssq (- n 1))))) Why does (ssq 4000000) run out of memory? Roce (USNA) Star-Class 6 Fal 201 8/11

Stack space in recursive calls
This function does the same thing, but takes an *extra argument* that
serves as an accumulator.
;; Sum of squares using tail recursion
(define (ssq-better n accum)
 (if (= n 0)
 accum
 (ssq-better (- n 1)
 (+ (sqr n) accum))))
Now (ssq-better 4000000 0) actually works!

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Tail recursion The second version worked because there was no need to make a stack of recursive calls. A function is *tail recursive* if its output expression in every recursive case is only the recursive call. In Scheme, this means the recursive call is **outermost** in the returned expression. ssq-better is better because it is tail recursive!

Tail recursion for Fibonacci

To implement tail recursion we usually make a helper function:

```
(define (fib-tail-helper n i fib-of-i fib-of-i+1)
 (if (= i n)
    fib-of-i
    (fib-tail-helper
    n
    (+ i 1)
    fib-of-i+1
    (+ fib-of-i
        fib-of-i+1))))
The main function then becomes:
(define (fib-tail n) (fib-tail-helper n 0 1 1))
```

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