

## SLR(1)

SLR(1) parsers handle conflicts by using one token of look-ahead:

- If the next token is an outgoing edge label of that state, shift and move on.
- If the next token is in the *follow set* of a non-terminal *that we can reduce to*, then do that reduction.

Of course, there may still be conflicts, in which case the grammar is not SLR(1). More look-ahead may be needed.

LALR parsers are similar, but they use more specialized FOLLOW sets rather than the "global" follow sets that we have seen.

SI413 - Class 12

Fall 2011 2 / 8

## Parse Trees

Beefed-up calculator language

$\mathit{run}  ightarrow \mathit{ares}$ <code>STOP</code> $\mathit{run} \mid \mathit{ares}$ <code>STOP</code>	
ares $ ightarrow$ VAR ASN bres $\mid$ bres	
bres $ ightarrow$ bres <code>BOP</code> res $\mid$ res	
$\mathit{res}  ightarrow \mathit{res}$ COMP $\mathit{exp} \mid \mathit{exp}$	
exp $ ightarrow$ exp <code>OPA</code> term $\mid$ term	
term $ ightarrow$ term <code>OPM</code> factor $\mid$ factor	
factor $ ightarrow$ NUM $\mid$ VAR $\mid$ LP bres RP	

Download today's tarball and run make to get a parse tree for some string in this language.

We notice that the parse tree is large and unwieldy with many unnecessary nodes.

SI413 - Class 12

Abstract Syntax Tree Consider the program x := (5 + 3) \* 2; x - 7;. What should the AST for this look like? SI413 - Class 12 Fall 2011 4 / 8 Roche (USNA) **AST** Properties Remember, ASTs are not about the syntax! They remove syntactic details from the program, leaving only the semantics. Typically, we show ordering (e.g. of *ares*'s in the previous example) by nesting: the last child of a statement is the next statement, or null. Are ASTs language independent? Roche (USNA) SI413 - Class 12 Fall 2011 5 / 8 Static type checking Consider the string (7 > 2) + 3; This is an error. But where should this error be identified? Each node in the AST has a type, possibly "void". Roche (USNA) SI413 - Class 12 Fall 2011 6 / 8

## Static type checking with variables

What about the string x = 6 > 3; x \* 12;?

We have to know the *type* of the variable *x*. Otherwise, there is no way to detect this error at compile-time.

Only *statically-typed languages* allow this sort of checking. Remember, in this class *errors are a good thing!* 

Roche (USNA)

SI413 - Class 12

Fall 2011 7 / 8

## Class outcomes

You should know:

- What an AST is, and why we need them.
- The relationship between language, parse tree, and AST.
- How static type-checking works, at a basic level.

You should be able to:

- Draw a parse tree for a given string, given the grammar.
- Determine the AST from the parse tree. Note that there is some flexibility here!

SI413 - Class 12

Fall 2011 8 / 8