Class 10: Shift-reduce Parsing and CFSMs SI 413 - Programming Languages and Implementation

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Bottom-up Parsing

A bottom-up (LR) parser reads tokens from left to right and maintains a stack of terminal *and* non-terminal symbols.

At each step it does one of two things:

- Shift: Read in the next token and push it onto the stack
- **Reduce**: Recognize that the top of the stack is the r.h.s. of a production rule, and replace that r.h.s. by the l.h.s., which will be a non-terminal symbol.

The question is how to *build* an LR parser that applies these rules *systematically, deterministically,* and of course *quickly*.

Simple grammar for LR parsing

Consider the following example grammar:

 $S \rightarrow E$ $E \rightarrow E + T$ $E \rightarrow T$ $T \rightarrow n$

Examine a bottom-up parse for the string n + n.

How can we model the "state" of the parser?

Parser states

At any point during parsing, we are trying to expand one or more production rules.

The state of a given (potential) expansion is represented by an "LR item". For our example grammar we have the following LR items:

The • represents "where we are" in expanding that production.

Pieces of the CFSM

The CSFM (Characteristic Finite State Machine) is a FA representing the *transitions* between the LR item "states".

There are two types of transitions:

 Shift: consume a terminal or non-terminal symbol and move the • to the right by one.

Example:
$$T \rightarrow \bullet n$$
 $T \rightarrow n \bullet$

 Closure: If the • is to the left of a non-terminal, we have an ε-transition to any production of that non-terminal with the • all the way to the left.

Example:
$$E \rightarrow E + \bullet T$$
 $\stackrel{\epsilon}{\longrightarrow}$ $T \rightarrow \bullet n$

Nondeterministic CFSM for example grammar

